PSO 2

(Recursion Tree) Give a big-O closed form for each of the following recurrences. (Assume that T(x) = 1 for any $x \leq 1$.)

(1) $T(n) = 2T(n/4) + \sqrt{n}$

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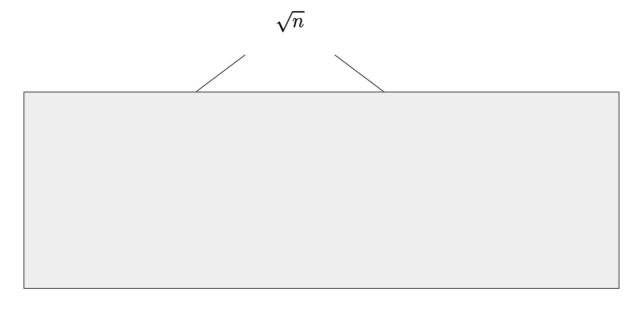
(2) T(n) = T(n/2) + T(n/3) + T(n/6) + n

Warning: Solving this T(n) using iterations is a bad idea!

... kind of, we will see that trees help us organize better!

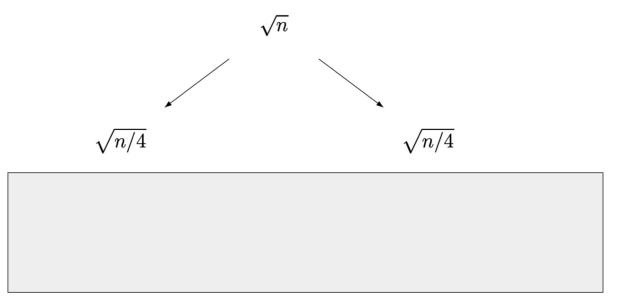
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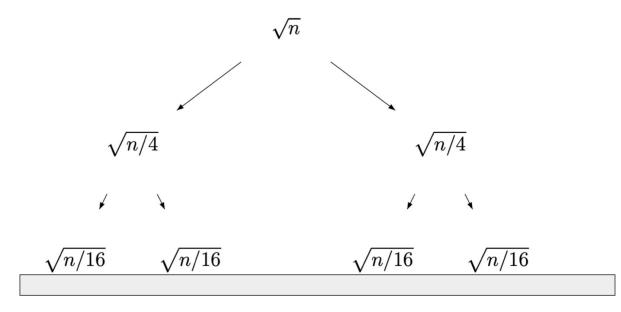
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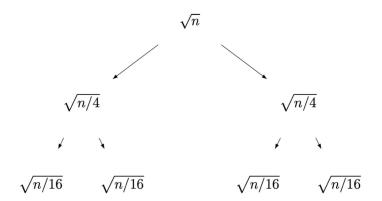
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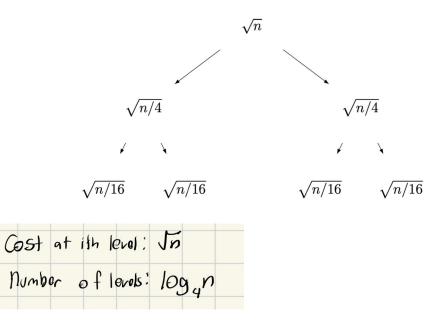


Cost at first level: Cost at second level: Cost at ith level: 1. Draw out the tree

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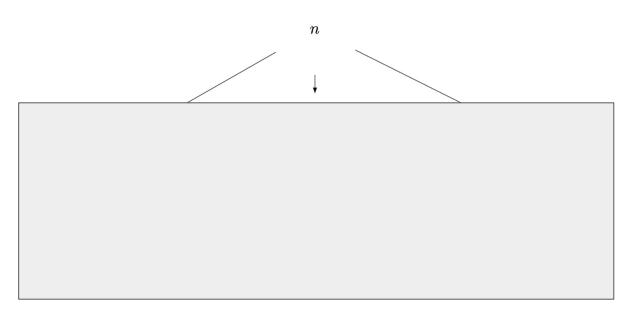
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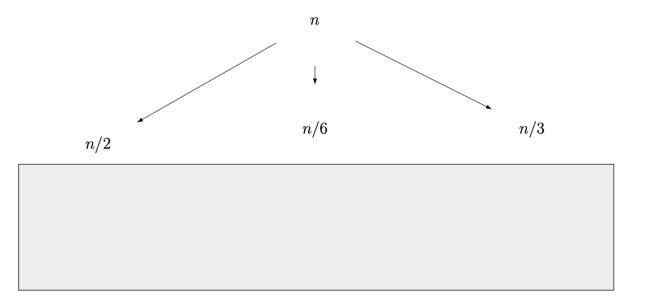
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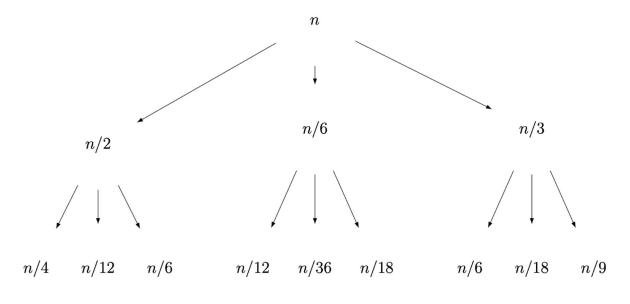
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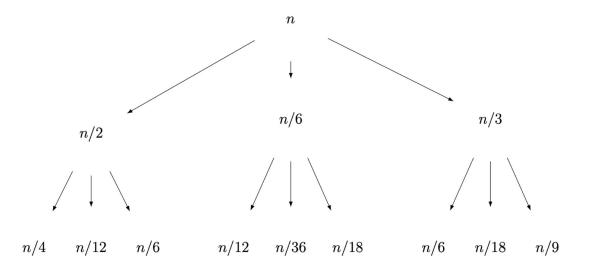
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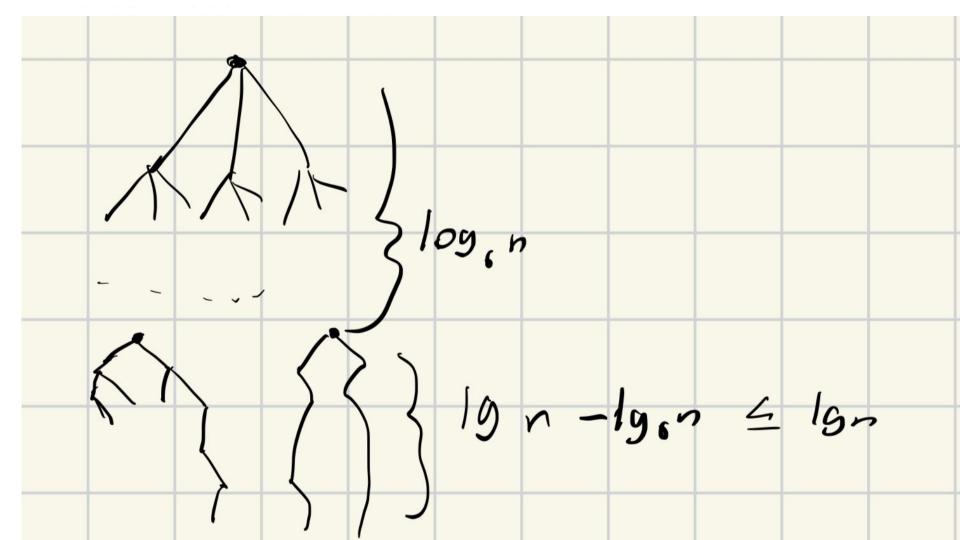


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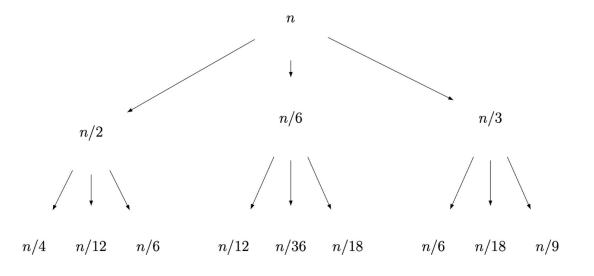
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Cost at ith level: N of levels: O(lgn) Number

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What is the problem with a tree?

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We usually like recurrences of this form

 $S(n) = \alpha S(n/\beta) + f(n),$

E.g question 1 recurrences

Solution: Variable change! But to what value?

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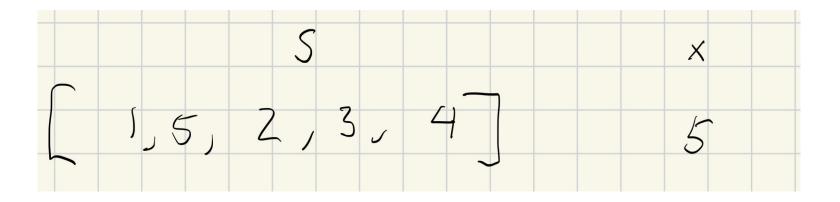
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This is just merge sort! O(mlogm) = O(log n * (log log n))

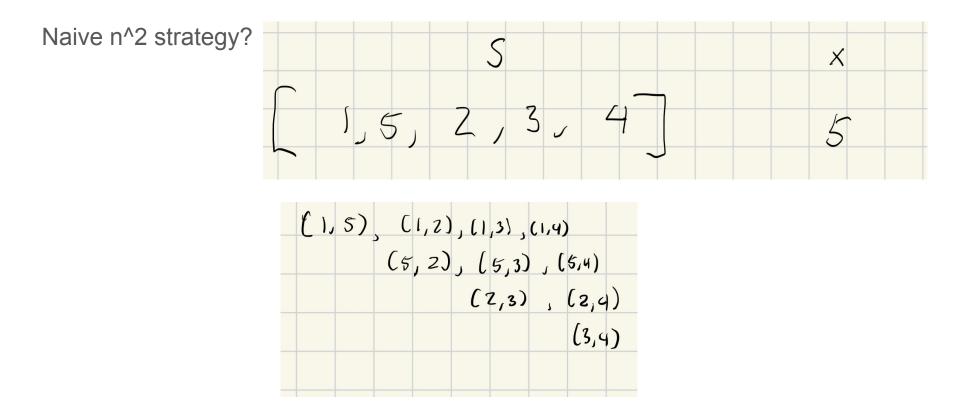
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Naive n^2 strategy?

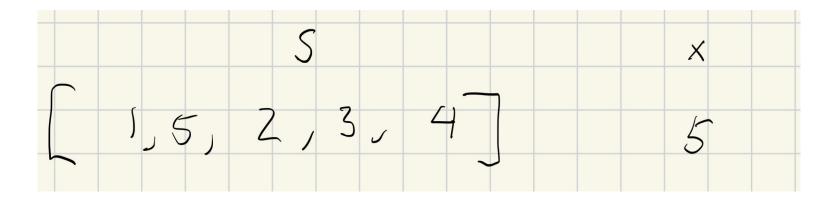


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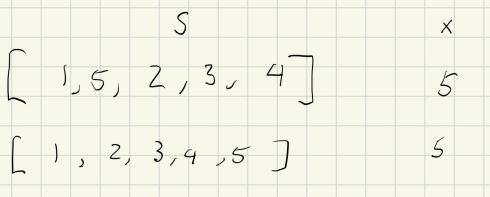


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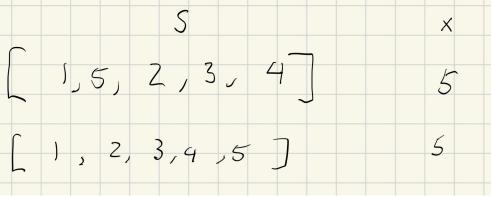


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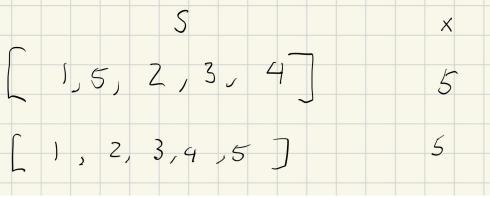
Idea: find pairs smarter

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Idea: find pairs smarter

Hold a left, right pointer, calculate sum

If sum > x: move right pointer

If sum < x: move left pointer

(Linked List) Consider a sorted circular doubly linked list of N numbers where the head element points to the smallest element in the list. Provide the asymptotic complexity in big- Θ with a brief explanation (including assumptions and analysis for each case, if there is more than one) for the following operations.

- 1. Inserting an element in its sorted position.
- 2. Finding the smallest element in the list.
- 3. Finding the 3^{rd} largest element in the list.
- 4. Finding the median in the list.